# CM0845 Logic Propositional Logic: Satisfiability

Andrés Sicard-Ramírez

Universidad EAFIT

Semester 2016-1

### Propositional Logic: Satisfiability

#### Remark

The reference for this section is Ben-Ari [2012, § 2.5].

# Satisfiability, Validity, Unsatisfiability and Falsifiability

Let  $\varphi \in PROP$ .

#### **Definitions**

- (i)  $\varphi$  is **satisfiable** iff  $[\![\varphi]\!]_v=1$  for some interpretation v. In this case, v is called a model for  $\varphi$ .
- (ii)  $\varphi$  is **valid** (a tautology), denoted  $\models \varphi$ , iff  $\llbracket \varphi \rrbracket_v = 1$  for all interpretations v.
- (iii)  $\varphi$  is **unsatisfiable** iff it is not satisfiable, that is, if  $[\![\varphi]\!]_v = 0$  for all interpretations v.
- (iv)  $\varphi$  is **falsifiable**, denoted  $\not\models$ , iff it is not valid, that is, if  $[\![\varphi]\!]_v=0$  for some interpretation v.

# Satisfiability, Validity, Unsatisfiability and Falsifiability

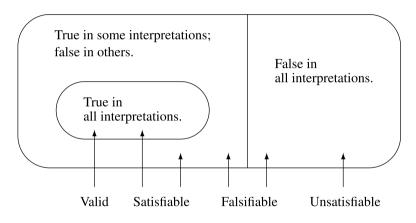


Figure 2.6 of [Ben-Ari 2012].

## Satisfiability, Validity, Unsatisfiability and Falsifiability

Theorem (Ben-Ari [2012], Theorem 2.39) Let  $\varphi \in PROP$ .

- (i) The proposition  $\varphi$  is valid if and only if  $\neg \varphi$  is unsatisfiable.
- (ii) The proposition  $\varphi$  is satisfiable if and only if  $\neg \varphi$  is falsifiable.

## Satisfiability of a Set of Propositions

Let  $\Gamma = \{\varphi_1, \dots\}$  be a set of propositions.

#### **Definitions**

- (i)  $\Gamma$  is **satisfiable** iff there exists an interpretation v such that  $[\![\varphi]\!]_v = 1$  for all  $\varphi_i \in \Gamma$ . In this case, v is a model of  $\Gamma$ .
- (ii)  $\Gamma$  is **unsatisfiable** iff for every interpretation v, there exists an  $\varphi_i \in \Gamma$  such that  $[\![\varphi]\!]_v = 0$ .

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#### Example

Prove that if  $\Gamma$  is unsatisfiable and for some i, the proposition  $\varphi_i$  is valid, then  $\Gamma - \{\varphi_i\}$  is unsatisfiable [Ben-Ari 2012, Exercise 2.15, p. 46].

#### References



Ben-Ari, Mordechai [1993] (2012). Mathematical Logic for Computer Science. 3rd ed. Springer (cit. on pp. 2, 4-7).