# CM0081 Formal Languages and Automata § 1.4 Formal Proofs

Andrés Sicard-Ramírez

Universidad EAFIT

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### Preliminaries

### Conventions

- The number and page numbers assigned to chapters, examples, exercises, figures, quotes, sections and theorems on these slides correspond to the numbers assigned in the textbook [Hopcroft, Motwani and Ullman (1979) 2007].
- The natural numbers include the zero, that is,  $\mathbb{N} = \{0, 1, 2, ...\}$ .

The power set of a set A, that is, the set of its subsets, is denoted by  $\mathcal{P}A$ .

### Proofs by Contradiction and Proofs of Negations

Proof by contradiction (or *reductio ad absurdum*)

$$\begin{bmatrix} \neg \beta \\ \vdots \\ \underline{\bot} \\ \beta \end{bmatrix}$$

Proof of negation [Bauer 2017]



### Proofs by Contradiction and Proofs of Negations

Proof by contradiction (or *reductio ad absurdum*)



Justifications



Proof of negation [Bauer 2017]



Proofs by Contradiction and Proofs of Negations

## Inductive Proofs: Mathematical Induction

### The induction principle

Let S(n) be a property about natural numbers. If

(i) we prove S(i) (basis step) and

(ii) we prove that for all natural number  $n \ge i$ , S(n) implies S(n+1) (inductive step),

then we may conclude S(n) for all  $n \ge i$ .

### The structural induction principle

Let  $S({\boldsymbol X})$  be a property about structures  ${\boldsymbol X}$  that are defined by some recursive/inductive definition. If

- (i) we prove S(X) for the basis structure(s) of X (basis step) and
- (ii) given a structure X whose recursive/inductive definition says it is formed from  $Y_1, \ldots, Y_k$ , we prove S(X) assuming that the properties  $S(Y_1), \ldots, S(Y_k)$  hold (inductive step),

then S(X) is true for all X.

#### Example

Given the functions  $f, g, h : \mathbb{N} \to \mathbb{N}$ ,

$$\begin{array}{ll} f(0)=0, & g(0)=1, & h(0)=0, \\ f(n+1)=g(n), & g(n+1)=f(n), & h(n+1)=1-h(n), \end{array}$$

and properties R, S, T,

 $R(n)\colon h(n)=1-g(n), \qquad \quad S(n)\colon h(n)=f(n), \qquad \quad T(n)\colon S(n)\wedge R(n),$ 

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- 1. To prove  $(\forall n)R(n)$  (impossible!)
- 2. To prove  $(\forall n)S(n)$  (impossible!)

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- 1. To prove  $(\forall n)R(n)$  (impossible!)
- 2. To prove  $(\forall n)S(n)$  (impossible!)
- 3. To prove  $(\forall n)T(n)$  (by mutual induction)

Inductive Proofs

Proof



Proof

- **b** Basis step T(0). Easy.
- lnduction step  $T(n) \Rightarrow T(n+1)$ :

$$\begin{split} S(n): \quad h(n+1) &= 1-h(n) & \qquad (\text{def. of } h) \\ &= g(n) & \qquad (\text{IH } R(n)) \\ &= f(n+1) & \qquad (\text{def. of } f) \end{split}$$

$$\begin{split} R(n): \quad h(n+1) &= 1 - h(n) & (\text{def. of } h) \\ &= 1 - f(n) & (\text{IH } S(n)) \\ &= 1 - g(n+1) & (\text{def. of } g) \end{split}$$

### References

Bauer, A. (2017). Five States of Accepting Constructive Mathematics. Bulletin of the American Mathematical Society 54.3, pp. 481–498. DOI: 10.1090/bull/1556 (cit. on pp. 3, 4).



Hopcroft, J. E., Motwani, R. and Ullman, J. D. [1979] (2007). Introduction to Automata Theory, Languages, and Computation. 3rd ed. Pearson Education (cit. on p. 2).