CM0246 Discrete Structures Lattices

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Preliminaries

Convention

The number assigned to chapters, examples, exercises, figures, sections, and theorems on these slides correspond to the numbers assigned in the textbook (Rosen 2004).

Definition

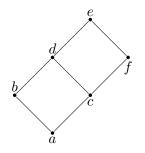
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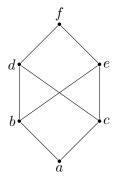
Example

The following poset is a lattice.



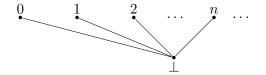
Example (counter-example)

The following poset is not a lattice because the upper bounds of the pair $\{b, c\}$ are d, e and f, but this set has not a least upper bound.



Example (counter-example)

The following poset is not a lattice because for example, the pair $\{1,2\}$ has not supremum.



Example

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- Let A be a set. Is $(P(A), \subseteq)$ a lattice?

Algebraic Structures

Definition

An algebraic structure on a set $A \neq 0$ is essentially a collection of *n*-ary operations on A (Cohn 1981, p. 41).

Example

A semigroup (S, *) is a set S with an associative binary operation $* : S \times S \to S$.

Example

A monoid $(M, *, \epsilon)$ is a semigroup (M, *) with an element $\epsilon \in M$ which is an unit for *, i.e. $\forall x(x * \epsilon = \epsilon * x = x)$.

Definition

Let \wedge and \vee be two binaries operations, called meet and join, respectively. A **lattice** retículo is an algebraic structure (L, \wedge, \vee) , which satisfy the following axioms for all x, y and z in L (Lipschutz and Lipson 2007):

$x \wedge y = y \wedge x$	(Commutative laws)
$x \lor y = y \lor x$	
$(x \land y) \land z = x \land (y \land z)$	(Associative laws)
$(x \lor y) \lor z = x \lor (y \lor z)$ $x \land (x \lor y) = x$	(Absortion laws)
$\begin{array}{l} x \land (x \land y) = x \\ x \lor (x \land y) = x \end{array}$	

Example

Let A be a set. $(P(A), \cap, \cup)$ is a lattice.

Definition

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 is $x \lor (y \land x) = x \land x$.

Theorem (principle of duality)

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Theorem (principle of duality)

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Proof.

The dual of every axiom in a lattice is also an axiom. Hence, the dual theorem can be proved by using the dual of each step of the proof of the original theorem.

Example

Let (L, \wedge, \vee) be a lattice. Prove the idempotent laws

$$x \wedge x = x, \tag{1}$$

$$x \vee x = x. \tag{2}$$

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$$\begin{aligned} x \wedge x &= x, \\ x \vee x &= x. \end{aligned}$$
 (1)

Proof of (1).

$$\begin{aligned} x \wedge x &= x \wedge (x \vee (x \wedge y)) \\ &= x \end{aligned}$$

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Proof of (2). By principle of duality on (1).

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Prove that if x and y are elements of a lattice (L, \wedge, \vee) then $x \vee y = y$, if and only if, $x \wedge y = x$.

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Let's suppose $x \lor y = y$. Then

$$\begin{aligned} x &= x \land (x \lor y) \\ &= x \land y \end{aligned}$$

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$$\begin{aligned} x &= x \land (x \lor y) & (\text{first absortion law}) \\ &= x \land y & (\text{hypothesis}) \end{aligned}$$

Continued on next slide

 $\mathsf{Proof} \leftarrow.$

Let's suppose $x \wedge y = x$. Then

$$y = y \lor (y \land x)$$

= $y \lor (x \land y)$
= $y \lor x$
= $x \lor y$

(second absortion law) (commutative law) (hypothesis) (commutative law) Theorem

Let (L, \wedge, \vee) be a lattice. Then (L, \preceq) is a partial order, where the relation \preceq is defined by (Lipschutz and Lipson 2007):

$$x \preceq y \stackrel{\mathsf{def}}{=} x \land y = x.$$

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Proof.

1. The relation \leq is reflexive

 $x \wedge x = x$ (idempotency), for all $x \in L$. Therefore $x \preceq x$, for all $x \in L$.

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Equivalence of the Definitions

Proof (continuation)

2. The relation \leq is antisymmetric

Suppose $x \leq y$ and $y \leq x$, then $x \wedge y = x$ and $y \wedge x = y$. Therefore

$x = x \wedge y$	(hypothesis)
$= y \wedge x$	(commutative law)
= y	(hypothesis)

That is, \leq is antisymmetric.

Continued on next slide

Equivalence of the Definitions

Proof (continuation).

3. The relation \leq is transitive

Suppose $x \leq y$ and $y \leq z$, then $x \wedge y = x$ and $y \wedge z = y$. Therefore

$x \wedge z = (x \wedge y) \wedge z$	(hypothesis)
$= x \wedge (y \wedge z)$	(associativity law)
$= x \wedge y$	(hypothesis)
= x	(hypothesis)

That is, $x \preceq z$.

Remark

Let (L, \wedge, \vee) be a lattice and let be (L, \preceq) the order partial induced by (L, \wedge, \vee) . It is possible prove that (L, \preceq) is a lattice.

Equivalence of the Definitions

Theorem (Problem 39, p. 500)

Let (L, \preceq) be a lattice. Then (L, \wedge, \vee) is a lattice, where

$$x \wedge y \stackrel{\text{def}}{=} \inf(x, y),$$
$$x \vee y \stackrel{\text{def}}{=} \sup(x, y),$$

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Proof.

1. Commutative laws for \land and \lor (Rosen's solution).

Because $\inf(x, y) = \inf(y, x)$ and $\sup(x, y) = \sup(y, x)$, it follows that $x \wedge y = y \wedge x$ and $x \vee y = y \vee x$.

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Proof (continuation)

2. Associative laws for \land and \lor (Rosen's solution).

Using the definition, $(x \wedge y) \wedge z$ is a lower bound of x, y and z that is greater than every other lower bound. Because x, y and z play interchangeable roles, $x \wedge (y \wedge z)$ is the same element.

Proof (continuation)

2. Associative laws for \land and \lor (Rosen's solution).

Using the definition, $(x \wedge y) \wedge z$ is a lower bound of x, y and z that is greater than every other lower bound. Because x, y and z play interchangeable roles, $x \wedge (y \wedge z)$ is the same element.

Similarly, $(x \lor y) \lor z$ is an upper bound of x, y and z that is less than every other upper bound. Because x, y and z play interchangeable roles, $x \lor (y \lor z)$ is the same element.

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Proof (continuation).

3. Absortion laws for \land and \lor (Rosen's solution).

To show that $x \land (x \lor y) = x$ it is sufficient to show that x is the greatest lower bound of x, and $x \lor y$. Note that x is a lower bound of x, and because $x \lor y$ is by definition greater than x, x is a lower bound for it as well. Therefore, x is a lower bound. But any lower bound of x has to be less than x, so x is the greatest lower bound.

The second statement is the dual of the first; we omit its proof.

References

Cohn, P. M. (1981). Universal Algebra. Revised edition. Vol. 6. Mathematics and Its Applications. D. Reidel Publishing Company (cit. on p. 9).



Lipschutz, S. and Lipson, M. L. (2007). Schaum's Outline of Discrete Mathematics. 3rd ed. McGraw-Hill (cit. on pp. 10, 12–15, 23, 24).



Rosen, K. H. (2004). Matemática Discreta y sus Aplicaciones. Trans. by Pérez Morales, J. M., Moro Carreño, J., Lías Quintero, A. I. and Ramos Alonc, P. A. 5th ed. McGraw-Hill (cit. on p. 2).